

Some Networking Aspects of Multiple Access

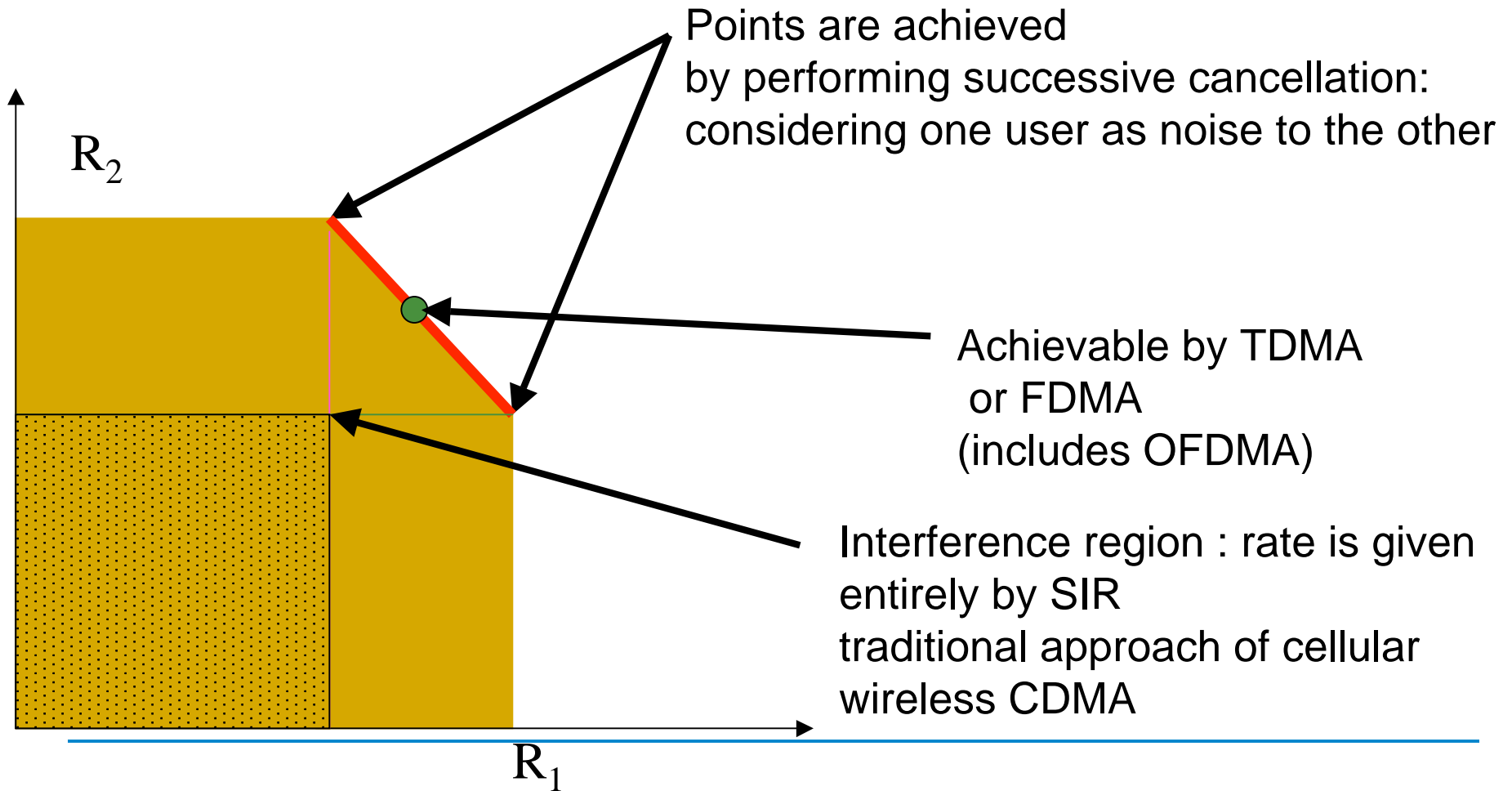
Muriel Medard

EECS

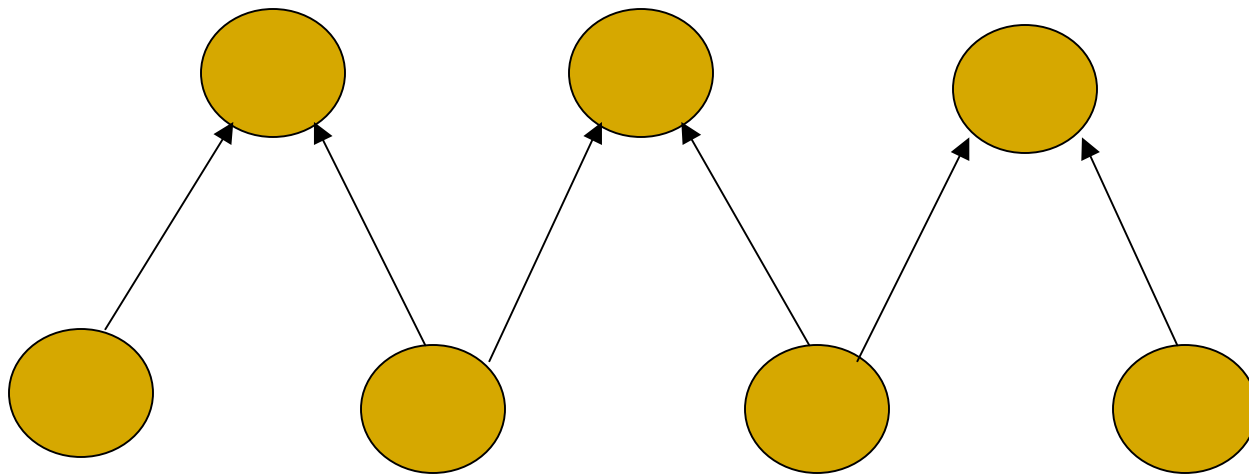
MIT

Multiple access

The Cover-Wyner rate region

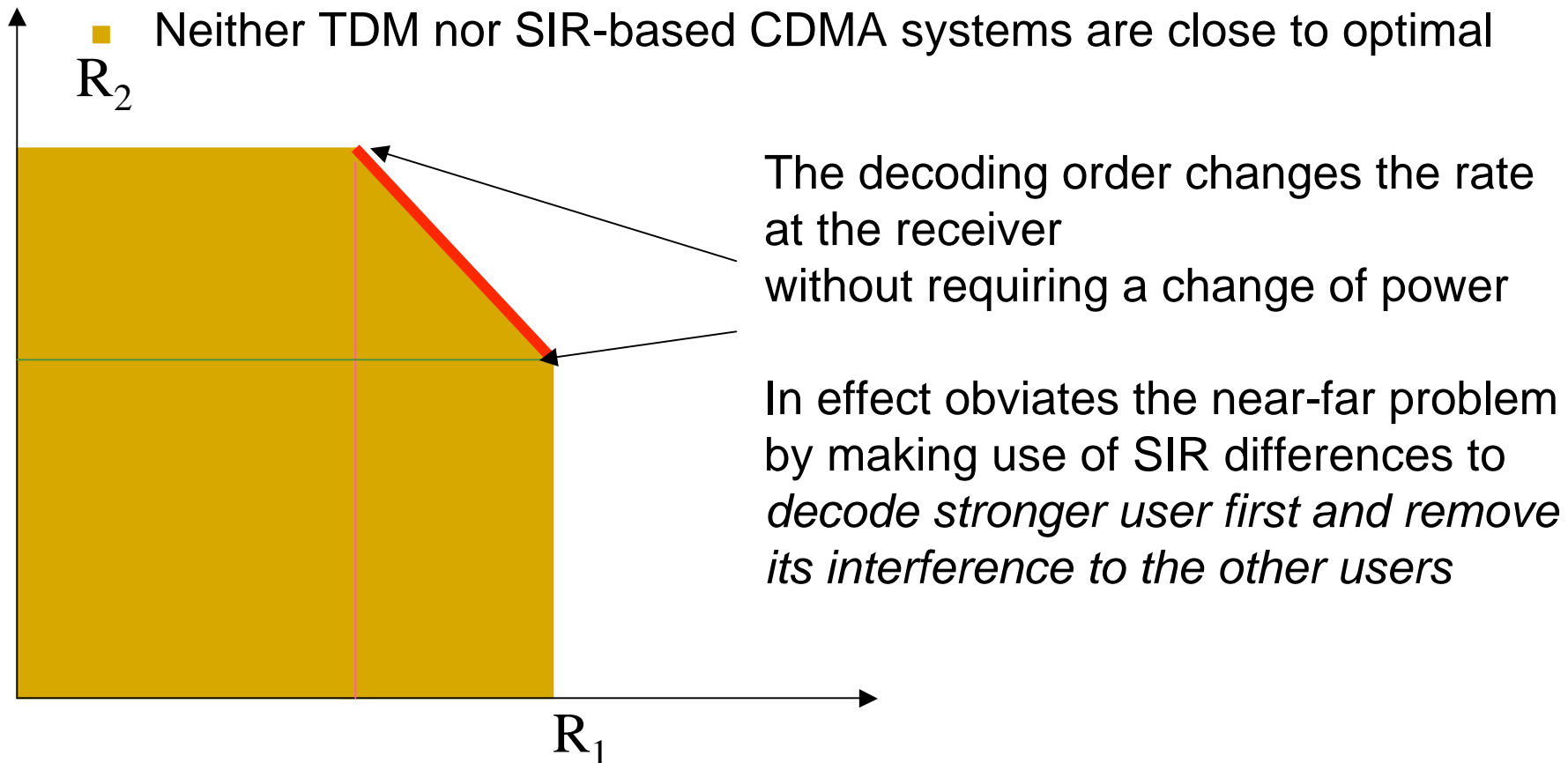


Multiple access in networks



Problem:
 Multiple access in uncertain conditions and channels
 Effect of multiple access on the network
 Interaction of several multiple access points

Intermediate SNRs



All rates on the maximum sum rate are achievable using rate-splitting

MAC Resource Allocation

- Utility maximization in non-fading channel
 - Gradient projection-like method with approximate projection
- Dynamic resource allocation in fading channel
 - Greedy vs. Optimal policy
 - Gradient projection method to track the greedy policy
- Two main approaches
 - Communications theory approach
 - No interference cancellation: CDMA [ODW'03], [KH'00]
 - TDMA [WG'05]
 - Information theoretic approach (non-concave utility)
 - Weighted sum rate maximization
 - D. Tse and S. Hanly, "Multi-Access Fading Channels: Part I: Polymatroid Structure, Optimal Resource Allocation and Throughput Capacities".
 - S. Vishwanath, S. A. Jafar, A. Goldsmith, "Optimum Power and Rate Allocation Strategies for Multiple Access Fading channels"

Non-fading channel

- The utility maximization problem

$$\text{maximize } u(\mathbf{R})$$

$$\text{subject to } \sum_{i \in S} R_i \leq \mathcal{C}\left(\sum_{i \in S} P_i, N_0\right), \quad \text{for all } S \subseteq \{1, \dots, M\}$$

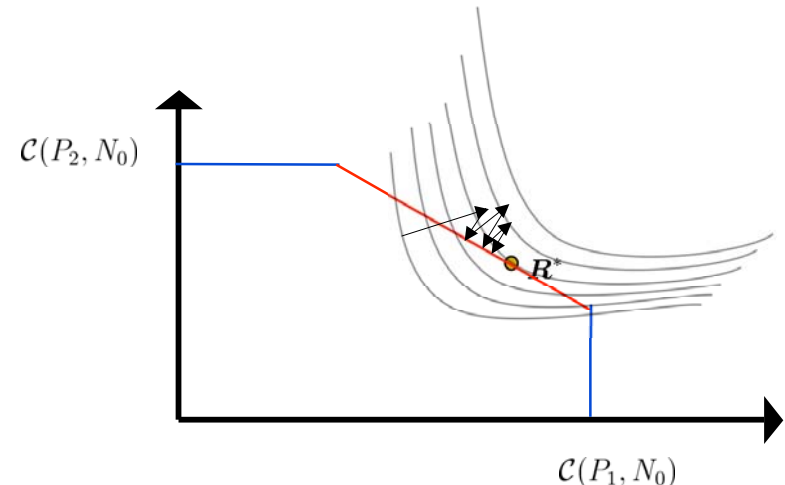
$$R_i \geq 0$$

- Assumption:** Let $u : \mathbb{R}^M \rightarrow \mathbb{R}$ be, monotonically non-decreasing and has bounded subgradient
- To solve this problem we use **gradient projection** method

$$\mathbf{R}^{k+1} = P(\mathbf{R}^k + \alpha^k \mathbf{g}^k)$$

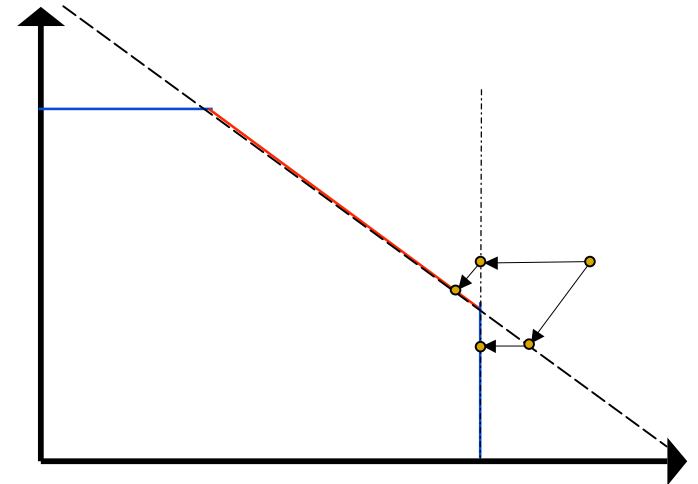
$$\mathbf{g}^k \in \partial u(\mathbf{R}^k)$$

- This method converges to the optimal solution for appropriate stepsize and bounded subgradient,



The Projection Problem

- Gradient projection effective when projection is easy
- In our problem, exact projection is hard
- **Main novelty: Approximate projection**
 - Successively project on the hyper planes corresponding to the violated constraints
 - Terminates after finite iterations
 - Resulting solution not unique

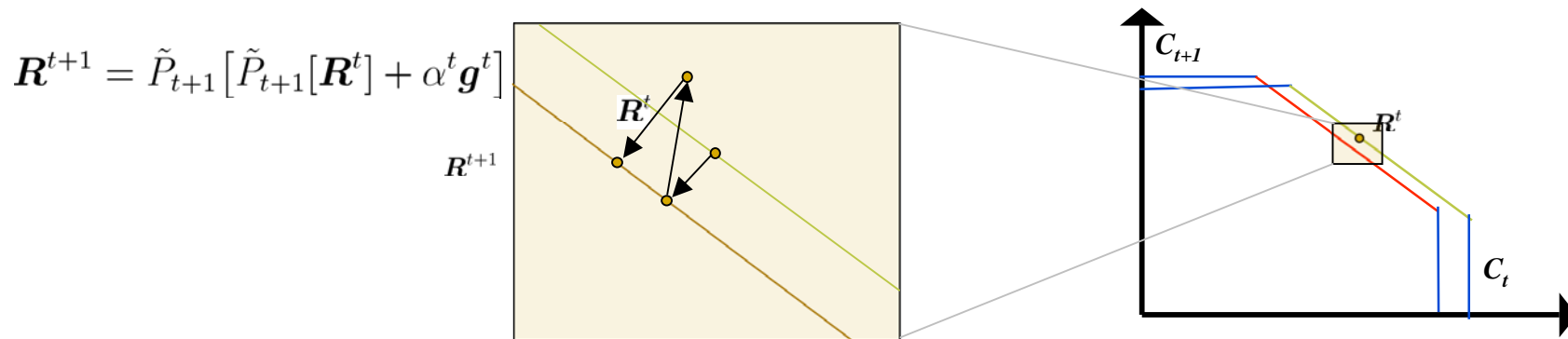


Fading Channel

- **Goal:** Find **rate allocation policy** $\mathcal{R} : \mathbb{R}^M \rightarrow \mathbb{R}$ s.t.

$$\mathbb{E}_{\mathbf{H}}[\mathcal{R}^*(\mathbf{H})] = \operatorname{argmax} u(\mathbf{R}) \quad \text{subject to } \mathbf{R} \in C_a(\mathbf{P})$$

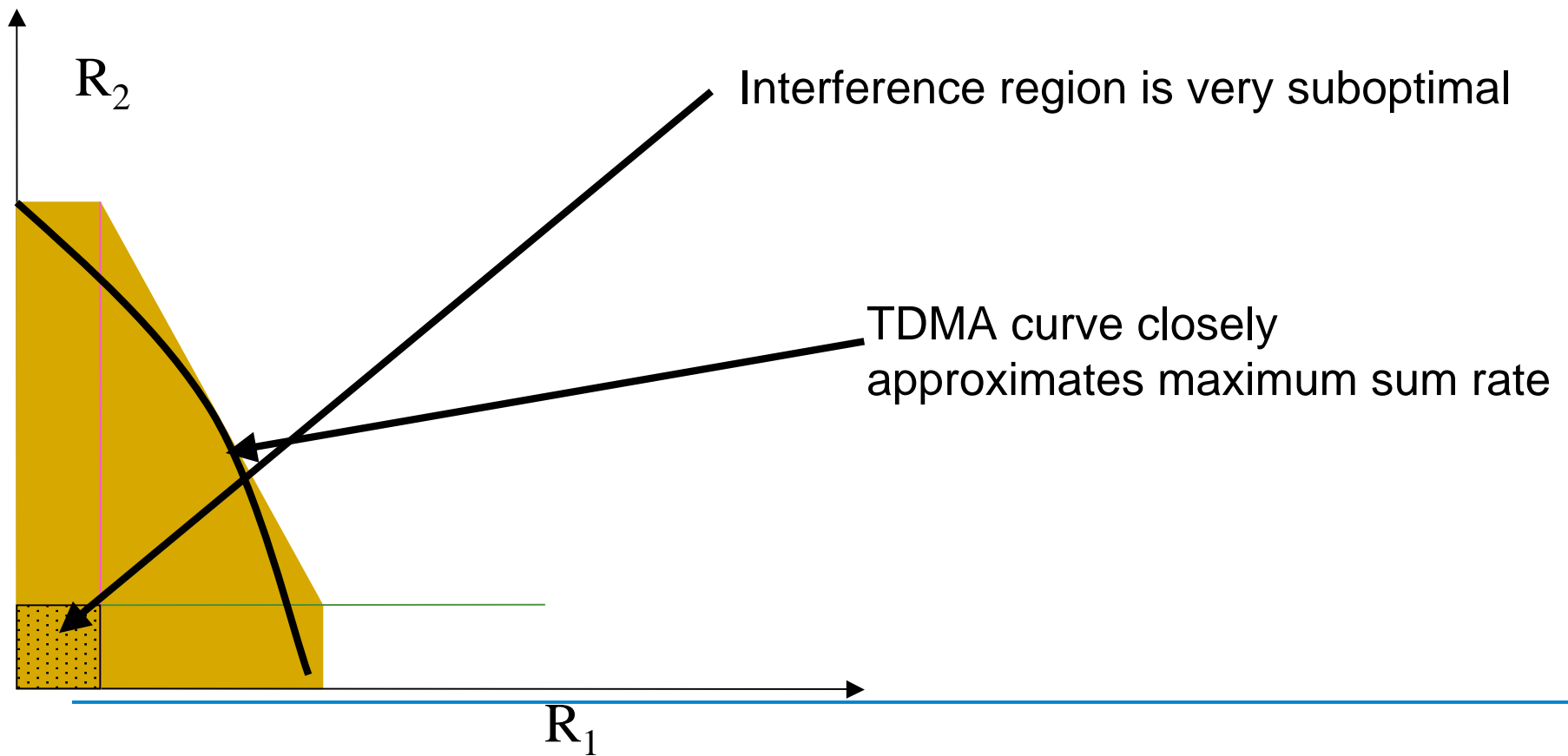
- There exists a **greedy policy** closely approximates the optimal policy.
- The greedy policy needs to solve the utility maximization problem at any time instant
 - ⇒ Not computationally efficient
- Under slow fading condition we can track the greedy policy by taking one iteration of the gradient projection method



- The rate allocation policy: $\tilde{\mathcal{R}}(\mathbf{H}(t)) = \mathbf{R}^t$

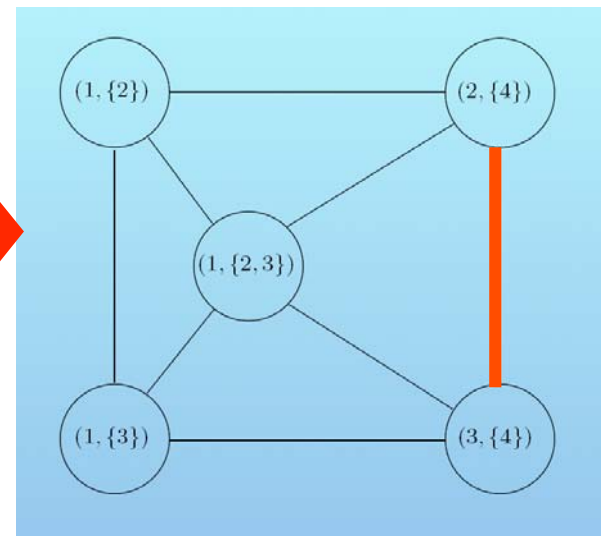
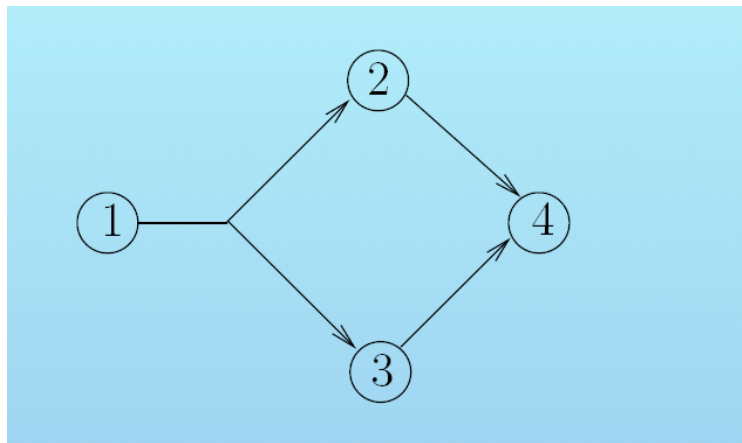
High SNR Case

In the case of high SNR, system is quasi-optimal if run as a TDMA system



Multiple access as a conflict

- A vertex represents a configuration of the network in terms of codes among components of the network
 - An edge between two vertices if the two configuration cause conflict (i.e. They cannot be served simultaneously).
- Stable Set: no edge connecting any pair of vertices in the set
 - Stable set of $G \rightarrow$ valid network configuration

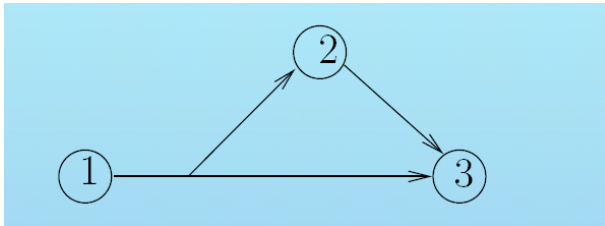


Traskov, Koetter, Medard

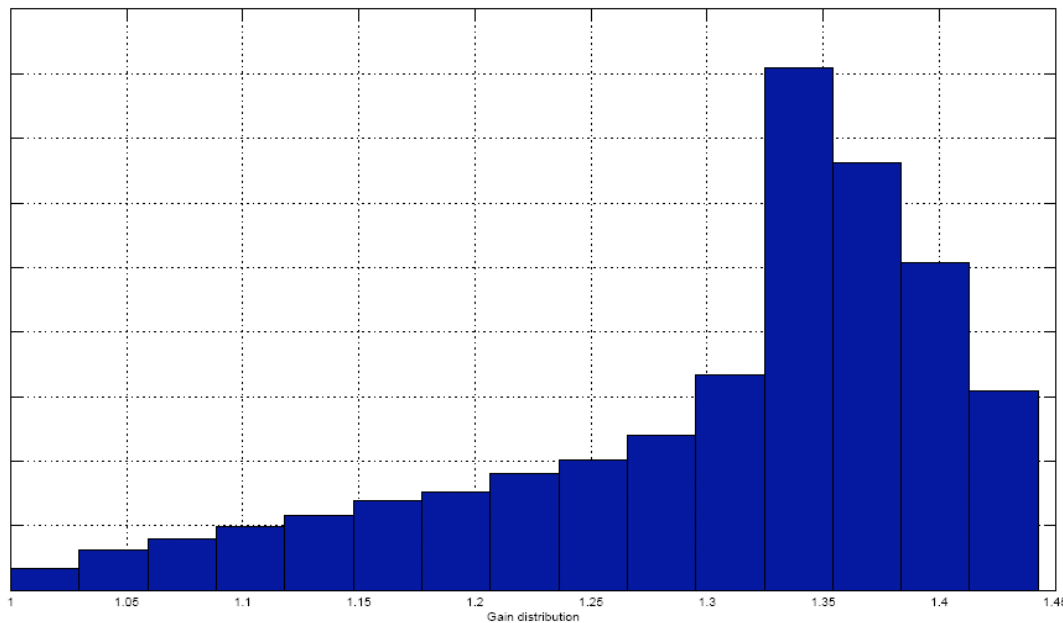
Multiple access as a conflict

- A valid transmission schedule is an *independent set* in the conflict graph.
 - The rate region is the convex hull of the *independent set polytope*.
 - In general this is computationally very demanding, however in certain graphs (perfect graphs) maximum independent sets can be computed efficiently (in polynomial time).
 - Examples of perfect graphs are *bipartite graphs* (e.g. layered networks), *fully connected graphs* (cluster networks), ...
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Network coding in a relay with conflicts

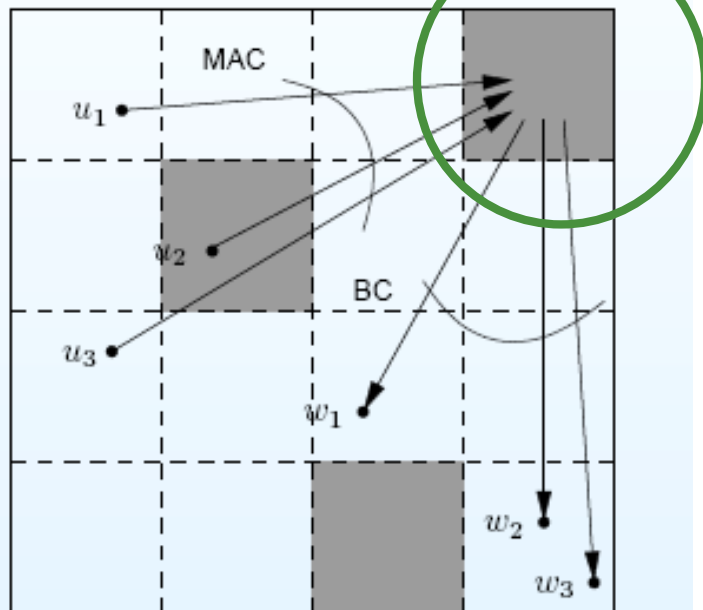


Take two nodes S and T at unit distance and place a relay randomly in the circle with midpoint $(S+T)/2$. We assume the erasures due to distance attenuation follow a square law rule. We compare the performance of network coding using the relay as opposed to routing.



Scaling laws using multiple access

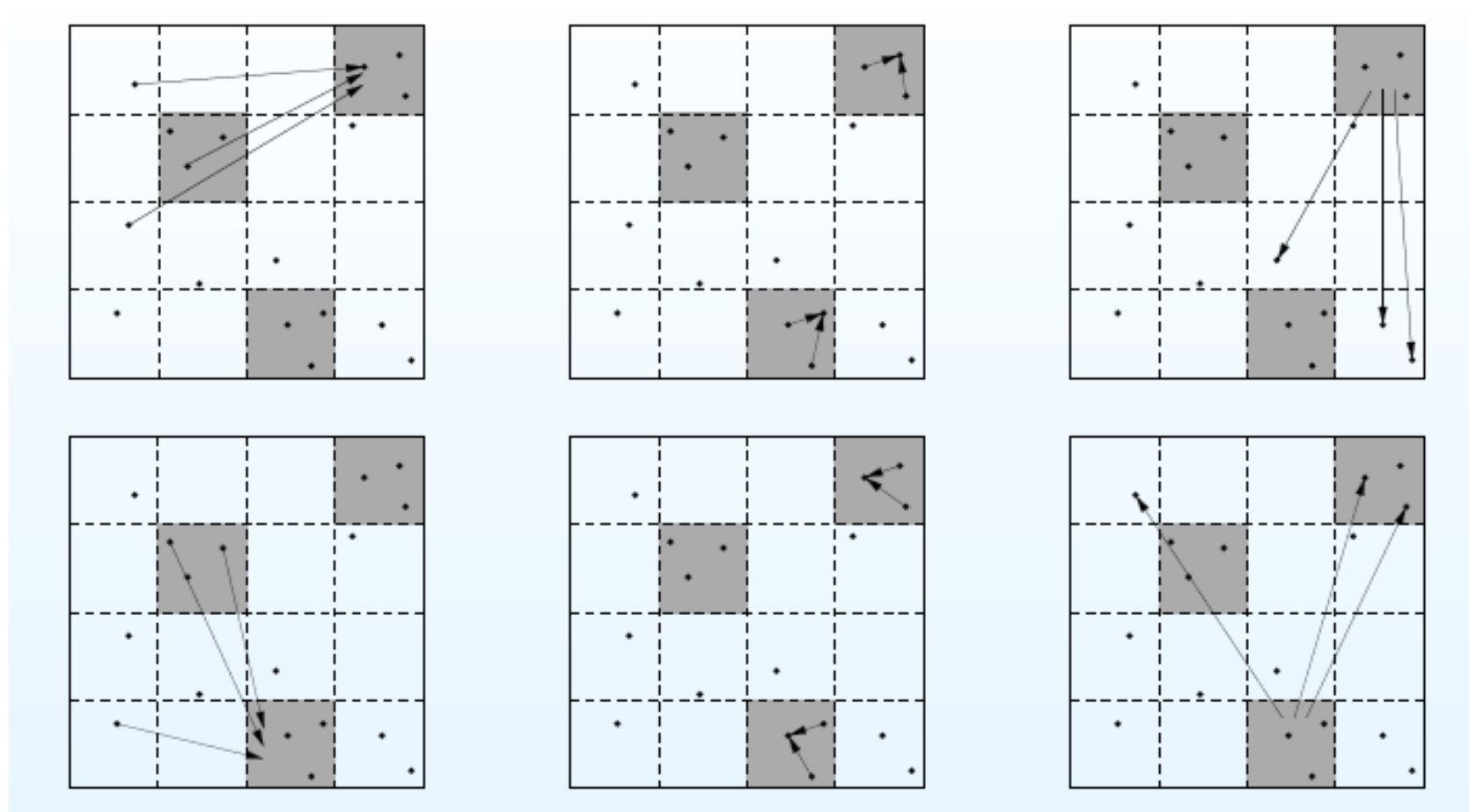
Squarelet acts as a component and interconnections occur among components, albeit in a multiple access fashion



- Example of granularity in systems
- Divide into squarelets
- Constant fraction of squarelets are dense (contain many nodes)
- Source-destination pairs relay traffic over dense squarelets
- Induces virtual multiple antenna multiple access and broadcast channels

Niessen, Shah, Gupta

Further hierarchical interaction



Consequences

- A hierarchical cooperative communication scheme, achieving a per node rate

$$\rho^{HR}(n) \geq n^{1 - \frac{\alpha}{2} - o(1)}$$

for any $\alpha > 2$, attenuation exponent

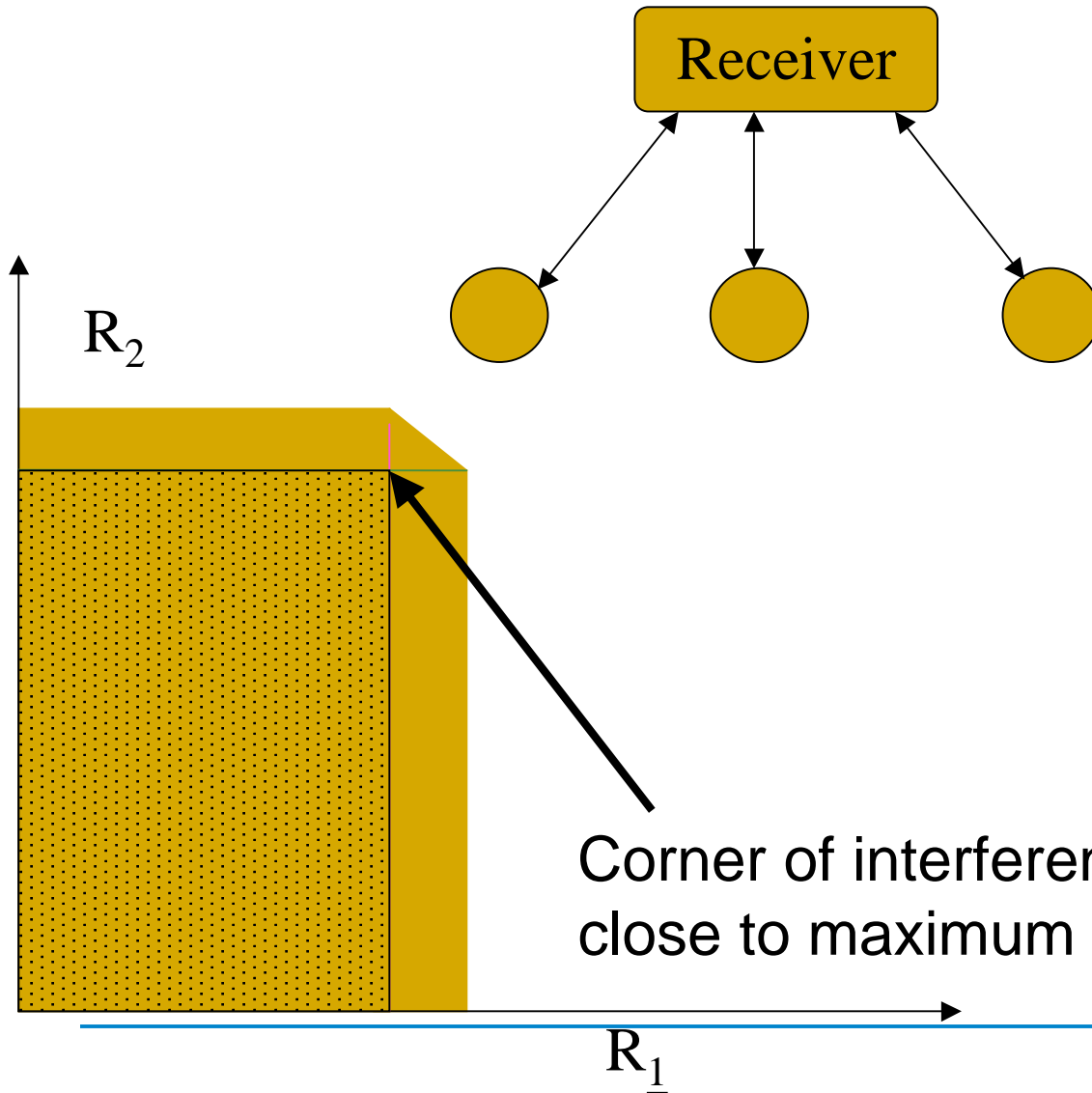
- For the best communication scheme

$$\rho^*(n) = O\left(n^{1 - \frac{\alpha}{2} + \varepsilon}\right)$$

for $\varepsilon > 0$ arbitrarily small and for any $2 < \alpha < 3$

- Thus scheme is order optimal for $2 < \alpha < 3$, so that decomposition is not detrimental in an order sense

Low SNR case - Interference



- The users compete for SIR at receiver
- The traditional cellular approach is a max-min approach based on explicit commands using closed-loop control
- No concept of priority, whether permanent or in response to a rapid change of circumstances

Corner of interference region is close to maximum sum rate

Going further

- Use of uncertain channel multiple access schemes in network settings
 - Placing soft constraints of multiple access in place of hard conflicts
 - Use of multiple access as an enabler for network coordination in moderate size networks
 - Low SNR networks – multiple access are locally obviated
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